### **BAPC 2019**

Solutions presentation

Problem Author: Ragnar Groot Koerkamp

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- The best team must play the most games, 2nd the 2nd-most, etc.
- For each node, find the longest path from each child to a leaf recursively.
- Extend the longest of these path to the current node.
- Sort both the list of path lengths and the skill levels.
- Match these lists one to one for a maximal solution.

Statistics: 15 submissions, 3 accepted, 8 unknown





Problem Author: Timon Knigge

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### B: Breaking Branches



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Given a branch of length n, determine who will win the game if they break it into pieces repeatedly.

- A branch of length n can be cut in exactly n-1 places.
- After any cut, the number of remaining possible cuts decreases by exactly one.
- lacksquare Alice wins when n is even. She can break it at any position.

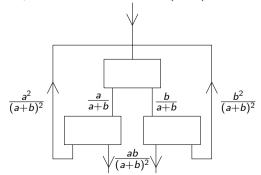
Statistics: 57 submissions, 56 accepted, 0 unknown

Problem Author: Daan van Gent

Given splitters that split their input producing an output ratio between their outputs of (a:b), can you build a network of splitters that produces an output ratio (c:d)?

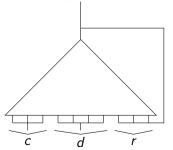
Problem Author: Daan van Gent

Given splitters that split their input producing an output ratio between their outputs of (a:b), can you build a network of splitters that produces an output ratio (c:d)? Step 1: Produce a ratio of (1:1).



Problem Author: Daan van Gent

Step 2: Create a binary tree of depth n of these (1:1) splitters, with n such that  $c+d \leq 2^n$ . Then connect c of the leafs to output 1, d leafs to output 2, and  $r=2^n-(c+d)$  back to the root.



Problem Author: Daan van Gent

Step 3: Remove all (1:1) splitters whose outputs are the same. To not run foul of timelimits, this needs to be done during generation of the tree.

Statistics: 1 submissions, 1 accepted, 0 unknown

#### D: Deck Randomisation

Problem Author: Ragnar Groot Koerkamp

Given two permutations A and B, how often do we need to repeat them one after the other to get back to where we started.

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Given two permutations A and B, how often do we need to repeat them one after the other to get back to where we started. Two possibilities:

- Option 1:  $(AB)^n = 1$ , minimum of 2n
- Option 2:  $(AB)^n A = 1$ , minimum of 2n + 1.

Both take similar strategy.

#### D: Deck Randomisation

Problem Author: Ragnar Groot Koerkamp

First, calculate the permutation AB, and split it into cycles. Example (notation from problem):

A = 516324

B = 465132

AB = 342561

Then AB has cycle (1 3 2 4 6 5)

#### For option 1:

- A cycle of length  $k_i$  implies  $k_i | n$ .
- Hence  $n = lcm(k_i)$ .
- **Example:** 1 cycle of length 6, so n = 6, giving 12 shuffles.

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#### For option 2:

- Per cycle of length  $k_i$ , check if exists z such that  $(AB)^z = A$ . Then  $n = -z \pmod{k_i}$ .
  - Example:  $(AB)^4 = 516324$ , so  $n = -4 \pmod{6}$ .
- Option 2 is not applicable if such z does not exist.
- Then reconstruct *n* using the Chinese Remainder Theorem.
  - **Example:** n = 2, giving 5 shuffles.

Minimum of the two is answer.

Statistics: 26 submissions, 1 accepted, 11 unknown

Problem Author: Raymond van Bommel



Compute the minimal number of coin exchanges needed to pay n.

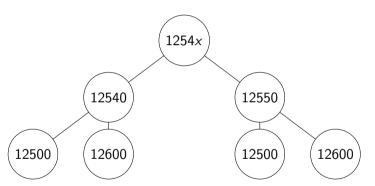
- Note that you never need to exchange more than 9 coins of value  $10^k$ , because for 10 exchanges we just use a single  $10^{k+1}$  coin.
- Consider 1254x. If we can do 12540 in a exchanges and 12550 in b, then we can do

$$\min\{a + x, b + 10 - x\}.$$

Solving recursively is costs 2<sup>1000</sup> calls.

## E: Efficient Exchange

Problem Author: Raymond van Bommel



There are only two distinct nodes per layer. So only 1000 calls needed.

Statistics: 124 submissions, 23 accepted, 37 unknown

Problem Author: Bjarki Ágúst Guðmundsson

• Given a family picture, find if there are three people of relative height 2, 1, and 3, in this order.

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- For each position, find the largest value on the right of it.
- Go from left to right and keep a set of all values seen so far.
- For each value, find the smallest element of the set that is larger.
- A 213 ordering exists if this smallest larger element on the left is smaller than the largest element on the right.

Statistics: 100 submissions, 20 accepted, 27 unknown

Problem Author: Mees de Vries

■ Your fungus is growing. How many squares does it occupy after *k* steps?

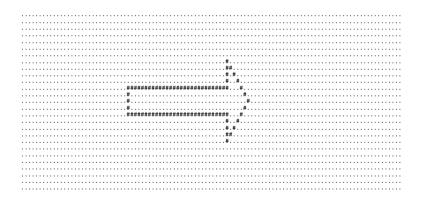
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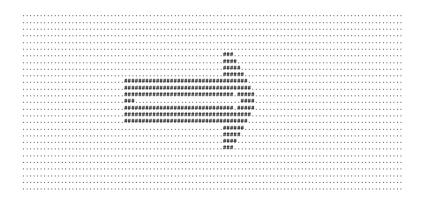
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- Your fungus is growing. How many squares does it occupy after *k* steps?
- Obvious solution: flood fill. But k can be  $10^6$ : too slow.
- We need to be smarter. Let's look at an example.

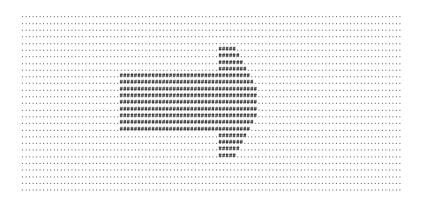
Problem Author: Mees de Vries



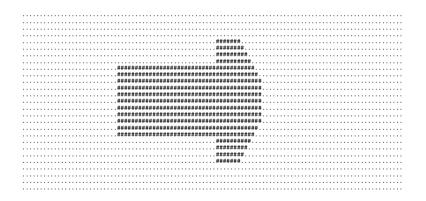
Problem Author: Mees de Vries



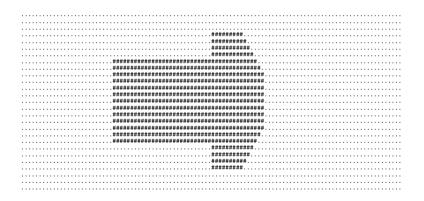
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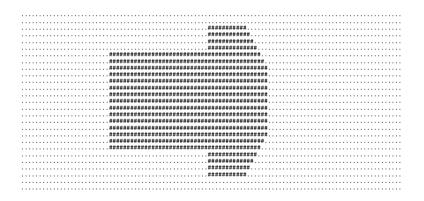
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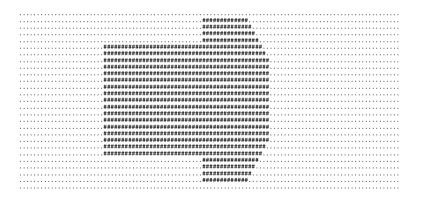
Problem Author: Mees de Vries



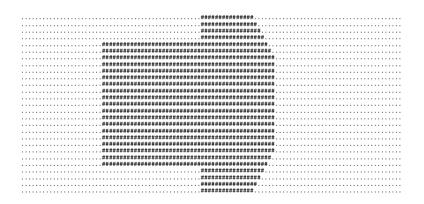
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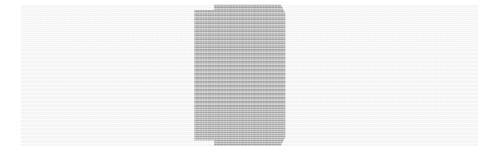


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Eventually everything turns into: a big rectangle with weird corners.



Problem Author: Mees de Vries

#### Solution:

- Simulate for 20 steps, then compute the height/width of the bounding rectangle.
- Count how many squares in the rectangle are missing.
- lacktriangle Find the height imes width for the final bounding rectangle. Subtract corner squares.

Statistics: 52 submissions, 10 accepted, 22 unknown

Problem Author: Bruno Ploumhans

Given k vases, and p pedestals, place every vase on a pedestal of the right size. Observations:

- If a pedestal has only one size, match those to vases of that size.
- Then you only have pedestals  $(1,2),(2,3),(3,4),\ldots$
- Match all 1-vases to (1, 2)-pedestals.
- Use leftover (1,2)-pedestals for 2-vases.
- Then you only have pedestals  $(2,3),(3,4),(4,5),\ldots$
- ... and repeat.

In other words: go greedily from left to right.

#### H: Historic Exhibition

Problem Author: Bruno Ploumhans

#### Common pitfalls:

- Output impossible once the smallest vase does not fit the smallest available pedestal. (There might still be larger pedestals!)
- Flow algorithm: too slow.

Statistics: 100 submissions, 26 accepted, 29 unknown

Problem Author: Timon Knigge

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- Solution:
  - Find the k additional edges, at most 16.
  - For each additional edge, at least one end point is not in the independent set. Brute force all  $2^k$  options.

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- Normally this is NP complete, but ... this graph is very close to a tree!
- We can solve this problem in linear time for a tree.
- Solution:
  - Find the k additional edges, at most 16.
  - For each additional edge, at least one end point is not in the independent set. Brute force all 2<sup>k</sup> options.
- Total runtime  $O(2^k \cdot n)$ .

Statistics: 20 submissions, 1 accepted, 7 unknown



Problem Author: Ragnar Groot Koerkamp

Given a squarefree number n, find an 1 < m < n such that  $n \times m$  is squarefree.



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- Squarefree test: trial division by all squares smaller than n.
- To find *m*: try all options starting at 2. This is fast since you'll allways find a solution among the first 13 primes.



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- Squarefree test: trial division by all squares smaller than n.
- To find *m*: try all options starting at 2. This is fast since you'll allways find a solution among the first 13 primes.
- Do **NOT** print n-1. It fails for  $k^2+1$ .

Statistics: 80 submissions, 56 accepted, 0 unknown



Problem Author: Ragnar Groot Koerkamp

```
static int primes[]={2,3,5,7,11,13,17,19,23,29,31,37, ..., 99991}
int main() {
    int n:
    cin >> n;
    for (int i = 0; i < (int)sizeof(primes) / (int)sizeof(primes[0]);</pre>
        i++) {
        if (n % primes[i] != 0) {
            cout << primes[i] << endl;</pre>
            break;
```



Problem Author: Ragnar Groot Koerkamp

```
t = int(input())
P = [2,3,5,7,11,13,17,19,23,29]
for i in P:
    if t%i!=0:
        print(i)
        break
```

print 3



Problem Author: Ragnar Groot Koerkamp

```
# Get input
n = int(raw_input())
# Find a number m such that m is relatively prime to n
# Check if any of these numbers divides n
if n %2 == 0:
    if n%3 == 0:
        if n\%5 == 0:
            if n\%7 == 0:
                if n %11 == 0:
                     if n %13 == 0:
                         if n %17 == 0:
                             print 19
                         else:
                             print 17
                     else:
                         print 13
                else:
                     print 11
            else:
                print 7
        else:
            print 5
    else:
```

### K: Keep Him Inside

Problem Author: Timon Knigge

- The problem was to find a weighted average of the vertices of some convex polygon that equals the given point *P*.
- Choose one of the vertices of the polygon as a "base point", translate so the base point is the origin and triangulate the polygon by drawing lines from the base point.
- Find the triangle in which the prisoner P lies (e.g. by calculating angles from the base point). Call the vectors from the base point to two other vertices of the triangle  $\mathbf{v}_1$  and  $\mathbf{v}_2$
- The vector P can now be decomposed into  $a \times \mathbf{v}_1 + b \times \mathbf{v}_2$  by projecting (calculate some inner products).
- The weights are 1 a b, a and b for the three points of the triangle and 0 otherwise.

Statistics: 53 submissions, 12 accepted, 28 unknown

Problem Author: Raymond van Bommel and Mees de Vries

At the casino, n players start with k lives each. Each round, each one loses a life with probability 1 - p. You win if you are the only one remaining. What is the probability of a draw?

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At the casino, n players start with k lives each. Each round, each one loses a life with probability 1-p. You win if you are the only one remaining. What is the probability of a draw?

$$\begin{split} \mathbb{P}(\mathsf{draw}) &= 1 - \mathbb{P}(\mathsf{someone\ wins}) \\ &= 1 - n \times \mathbb{P}(\mathsf{player\ 1\ wins}) \\ &= 1 - n \times \sum_{i=1}^{\infty} \mathbb{P}(\mathsf{player\ 1\ dies\ round\ } i, \ \mathsf{other\ player\ die\ before\ round\ } i) \\ &= 1 - n \times \sum_{i=1}^{\infty} \mathbb{P}(\mathsf{player\ 1\ dies\ round\ } i) \times \mathbb{P}(\mathsf{Player\ 1\ dies\ before\ round\ } i)^{n-1}. \end{split}$$

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For large i,  $\mathbb{P}(\text{player 1 dies round } i)$  is very small. So compute only for i up to M = 1000 (or more).

Problem Author: Raymond van Bommel and Mees de Vries

#### Two options:

Mathematically:

$$\mathbb{P}(\mathsf{Player}\ 1\ \mathsf{dies}\ \mathsf{round}\ i) = \binom{i-1}{k-1} p^{i-k} (1-p)^k.$$

This gives an  $\mathcal{O}(M)$  algorithm.

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This gives an  $\mathcal{O}(M)$  algorithm.

2 With dynamic programming: let

$$DP[r][I] = \mathbb{P}(Player 1 \text{ has } I \text{ lives in round } r.).$$

Then:

$$DP[r][I] = pDP[r-1][I] + (1-p)DP[r-1][I+1]$$

(plus correct edge conditions). This gives an  $\mathcal{O}(Mk)$  algorithm.

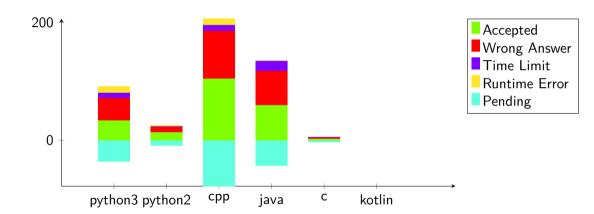
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Problem Author: Raymond van Bommel and Mees de Vries

Statistics: 3 submissions, 2 accepted, 0 unknown

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## Language stats



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#### The Proofreaders

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- Job Doesburg
- Nicky Gerritsen
- Raymond van Venetië
- Mees Vermeulen
- Jan Westerdiep

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- Daan van Gent
- Ragnar Groot Koerkamp
- Bjarki Ágúst Guðmundsson
- Joey Haas
- Timon Knigge
- Harry Smit
- David Venhoek
- Mees de Vries
- Wessel van Woerden