## KTH Challenge 2011 Solutions

#### Head of Jury: Lukáš Poláček Jury: Per Austrin, Emma Enström, Mikael Goldmann, Ulf Lundström

April 9, 2011

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- Just implement what the problem statement says.
- Can be made easier with bitset<32> in C++ or Integer.toBinaryString() in Java.

Statistics: 66 submissions, 39 correct, first at 00:03:29.

# E – Coast Length

• Surround the map from the input by empty squares. Flood it from one of the corners using a breadth-first-search.



- Time is linear in the size of the map.
- Watch out for stack size when implementing recursive Depth first search.

(4月) (日)

Statistics: 126 submissions, 17 correct, first at 0:49:43.

• The first 5 numbers are 1, 3, 5, 7, 9. Let's try with 11, 13, 15, ...

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## C – Base-2 Palindromes

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## C – Base-2 Palindromes

- The first 5 numbers are 1, 3, 5, 7, 9. Let's try with 11, 13, 15, ... Wrong Answer
- First find the bit length of the result. There are  $2^k$  palindromes of length 2k + 1 and 2k + 2.
- Now we have the length *L* and *M* the order of the palindrome among length-*L* palindromes. The left half of the palindrome is *M* in binary using *L* bits. Right half is a mirror of the left half.

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- First find the bit length of the result. There are  $2^k$  palindromes of length 2k + 1 and 2k + 2.
- Now we have the length L and M the order of the palindrome among length-L palindromes. The left half of the palindrome is M in binary using L bits. Right half is a mirror of the left half.

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#### • Accepted

Statistics: 63 submissions, 14 correct, first at 0:20:52.

- If a 1 is to the left of a 0, these two have to be swapped at some point. The same is true for 2/0 and 2/1.
- Process the sequence from left to right. Keep track of the number of 1's and 2's to the left of current number and calculate the result.
- Watch out for overflow.
- Linear time solution.

Statistics: 60 submissions, 12 correct, first at 1:13:57.

## G – Getting Rid of Coins

• Suppose the price is 100 and we have 50 1-cent coins? Would it make sense to use only 25 of them?

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- No. Since we will use at least one 5-, 10- or 25-cent coin, we would be better off paying with 1-cent coins instead of one big coin.

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- No. Since we will use at least one 5-, 10- or 25-cent coin, we would be better off paying with 1-cent coins instead of one big coin.
- We can generalize this argument for 5- and 10-cent coin.
- We try to pay using at least  $N_1 24$  1-cent coins; at least  $N_5 4$  5-cent coins,  $N_{10} 4$  10-cent coins and try to pay the rest with 25-cent coins.
- The total number of possibilities:  $25 \cdot 5 \cdot 5 = 625$ .
- Other solutions: dynamic programming and a greedy which tries only 2 possibilities.

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Statistics: 38 submissions, ?? correct, first at 0:51:06.

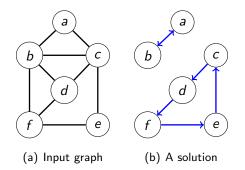
- We can assume that *s*<sub>1</sub> is in the first row. Rearranging columns doesn't change the result, so we can assume that the first row is sorted.
- Sorting items in the second row will not increase the area.
- Optimal solution will be always of the following form, for some k.

<i>s</i> <sub>1</sub>	 s <sub>k</sub>	$s_{2k+1}$	 <i>s</i> <sub>2<i>N</i>-1</sub>
$s_{k+1}$	 <b>s</b> <sub>2k</sub>	<i>s</i> <sub>2<i>k</i>+2</sub>	 s <sub>2N</sub>

• With some precomputation we can try each k in constant time.

Statistics: 16 submission, ?? correct, first at 2:54:54.

## B – Paintball

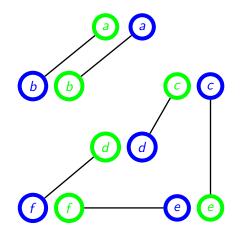


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 Instead of considering oriented edges, split each vertex into a green and blue vertex. Find a matching in the bipartite graph.



Statistics: 16 submissions, ?? correct, first at ??.

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## This was fun! When is the next contest?

- Programmeringstävlingsverksamhet trains every two weeks at KTH, check www.csc.kth.se/contest.
- Google Code Jam and TopCoder Open start in May. Both have usually around 10000 participants.
- Nordic Championships in October, Nort-western Europe qualifier in November.
- You will find more info at our webpage www.csc.kth.se/contest. Subscribe to our calendar and RSS feed.