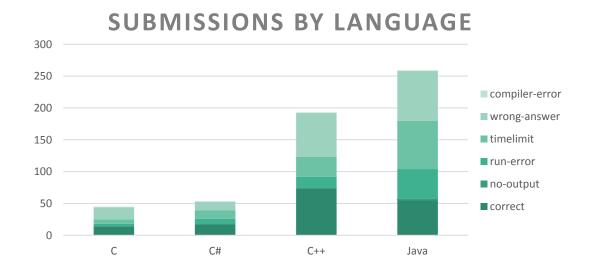
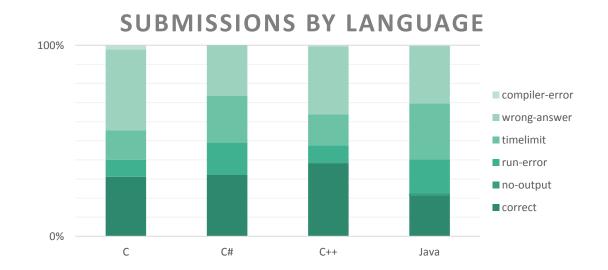
BAPC 2012

STATS + SOLUTIONS + SCORES

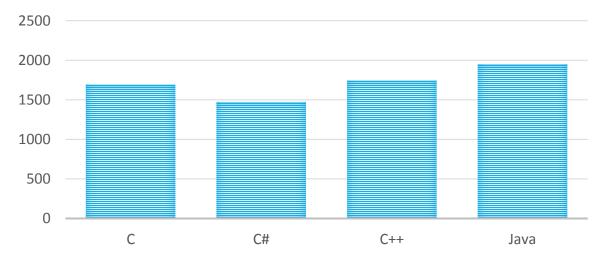
SUBMISSIONS OVER TIME



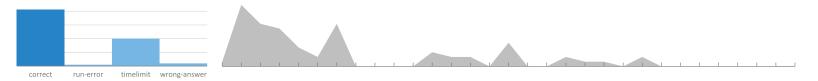




AVERAGE CODE LENGTH



Solutions



 \Rightarrow

Encoded Message

R	0	s	е	s
Α	r	U	R	е
d	V	i	0	1
е	t	S	Α	r
е	В	1	u	е

е	е	d	А	R
В	t	V	r	0
1	S	i	е	S
u	Α	0	R	е
е	r	1	е	S

int m = sqrt(message.length)

for i = 0 **to** m-1

for j = 0 **to** m-1

write(message[m-i-1+j*m])

writeline()

O(n)



Integer Lists

Actually reversing the lists is too slow

Use a **Deque** (double ended queue)

• Or keep 2 pointers to the begin and end yourself

O(n)



Good Coalition

Variant on Knapsack

Use dynamic programming

- dp(i, j) = maximum chance using exactly *j* seats considering only the first *i* parties
- Loop over parties and over number of seats to update this state
- Find maximum chance considering all parties for 76 .. 150 seats

 $O(n\cdot 150)$





Fire

First, do a Breadth First Search for the fires

- Start simultaneously from all fires
- For each floor, compute the time it will burn (distance to closest fire)

Then do another BFS to find the shortest path to the edge – without burning

Note that there might not be a fire

• And even then it might be impossible!

 $O(w \cdot h)$

correct run-error timelimit wrong-answer

Black Out

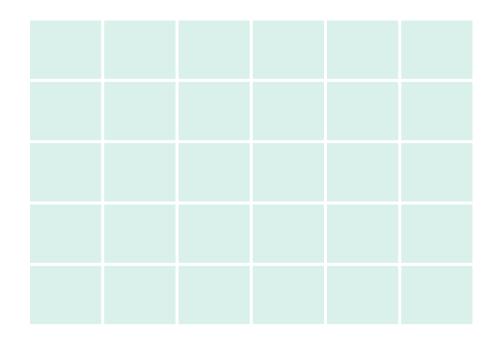
Greedy solution

- Use first move to black out one of the rows
- This makes the effective board dimensions even
- Now rotate every move of the opponent



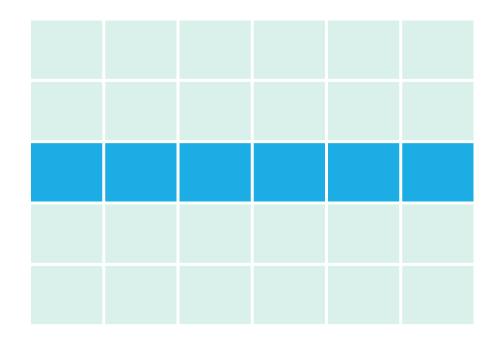






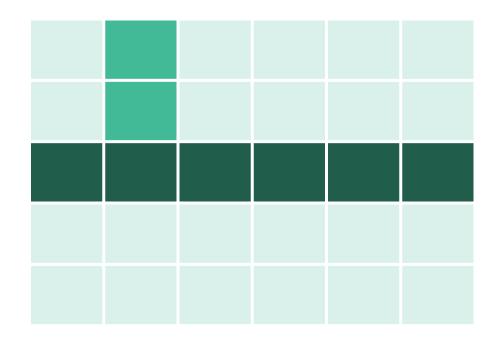






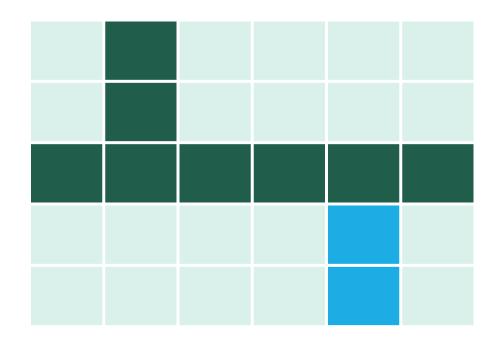






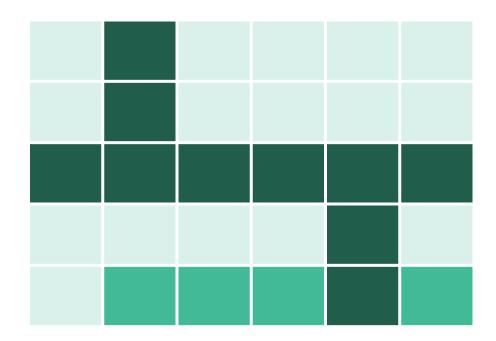






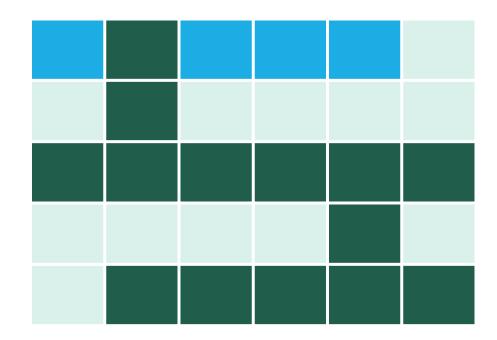






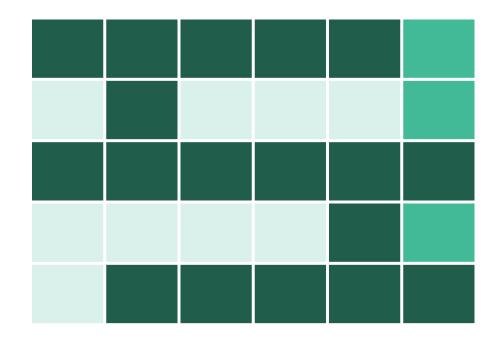






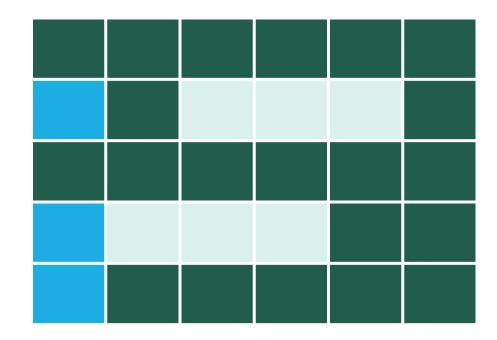


Black Out



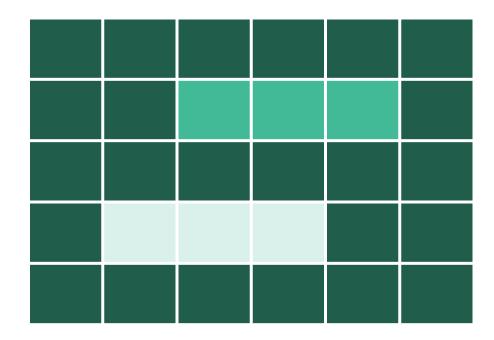


Black Out



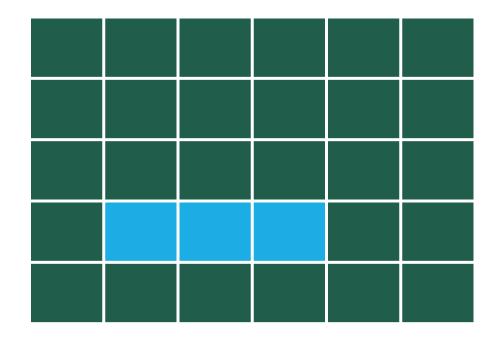


Black Out









correct run-error timelimit wrong-answer

Black Out

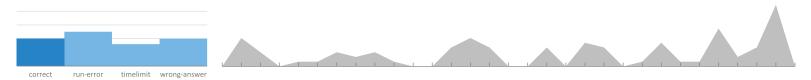
Greedy solution

- $\,\circ\,$ Use first move to black out one of the rows
- This makes the effective board dimensions even
- Now rotate every move of the opponent

Also possible to precompute entire state space

• Much harder to program

O(1) per move



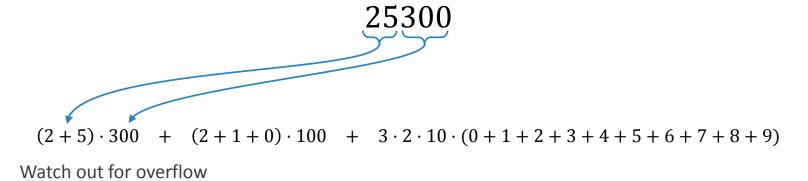
Digit Sum

O(b-a) is too slow!

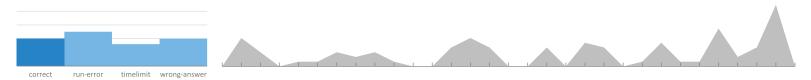
Simplify: instead of ds(a, b) do ds(0, b) - ds(0, a - 1)

Lower the digits to 0 one by one starting on the right

• For example to lower 25300 to 25000:



 $O(\log b)$



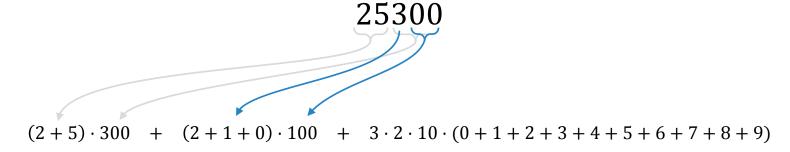
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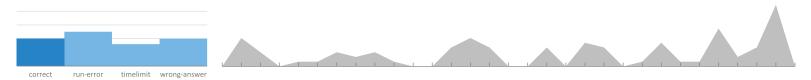
Lower the digits to 0 one by one starting on the right

• For example to lower 25300 to 25000:



Watch out for overflow

 $O(\log b)$



Digit Sum

O(b-a) is too slow!

Simplify: instead of ds(a, b) do ds(0, b) - ds(0, a - 1)

Lower the digits to 0 one by one starting on the right

• For example to lower 25300 to 25000:

 $(2+5) \cdot 300 + (2+1+0) \cdot 100 + 3 \cdot 2 \cdot 10 \cdot (0+1+2+3+4+5+6+7+8+9)$ Watch out for overflow

25300

 $O(\log b)$

prrect run-error timelimit wrong-answer

Another Dice Game



Another Dice Game

Either simulate all possible games with some tricks

- State: number of points, number of dice, subset of values left
- Use memoization to avoid redundant work
- $\,\circ\,$ With 8 dice only 1287 possible rolls (not enough time for $6^8 \approx 1.7 \text{M})$

 $O(40 \cdot 8 \cdot 2^{6} \cdot 1287)$

Or precompute team-side :)

- Only 40 inputs possible
- Even without smart dice rolls this takes less than 10 seconds





Chess Competition

For each player try to let him/her win

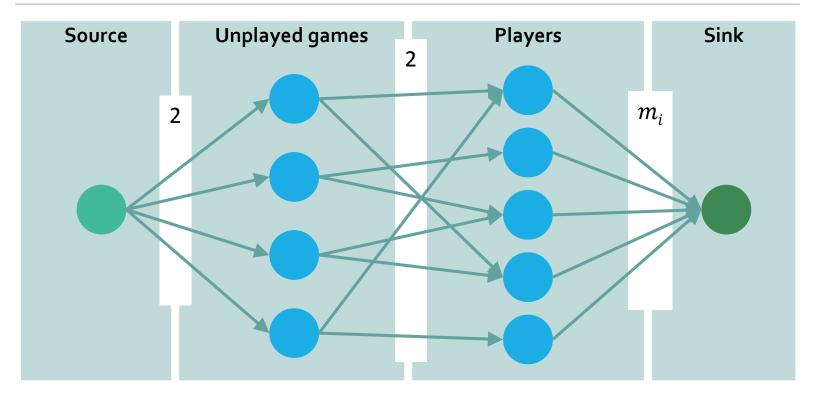
- Assume he/she wins all his/her remaining matches
- The points of the unplayed games should be distributed among other players
- But: in such way that nobody scores more points than this player
- Use Maximum Network Flow algorithm

 $O(n^5)$











John's Book Stack

Pulling out a book in the sorted top is never useful, so always pull out the first non-sorted book

However, simulation is too slow!

To move a big book from the top to position i, takes $2^i - 1$ steps

- (when there are no duplicates)
- Be a bit clever with duplicates

Watch out for overflow

O(n)

ect run-error timelimit wrong-answer

Hot Dogs in Manhattan

BFS to find minimum distance for each gridpoint

Binary search over answer *x*

Given an optimal solution, the leftmost point can be moved further to the left without loss of optimality if it doesn't move closer than *x* to another stand

- Same for the rightmost point to the right
- So: we only need to consider the leftmost and rightmost point with $\ge x$ per row

Now for each pair of rows, check for combination of leftmost and rightmost point

 $O(w \cdot h \cdot \log(w \cdot h))$







4	3	2	1	2	3	2	3	4
3	2	1		1	2	1	2	3
2	1	2	1	2	1		1	2
1		1	2	3	2	1	2	3
2	1	2	3	4	3	2	3	2
3	2	3	4	5	4	3	2	1
4	3	4	5	4	3	2	1	



4	3	2	1	2	3	2	3	4
3	2	1		1	2	1	2	3
2	1	2	1	2	1		1	2
1		1	2	3	2	1	2	3
2	1	2	3	4	3	2	3	2
3	2	3	4	5	4	3	2	1
4	3	4	5	4	3	2	1	



4	3	2	1	2	3	2	3	4
3	2	1		1	2	1	2	3
2	1	2	1	2	1		1	2
1		1	2	3	2	1	2	3
2	1	2	3	4	3	2	3	2
3	2	3	4	5	4	3	2	1
4	3	4	5	4	3	2	1	

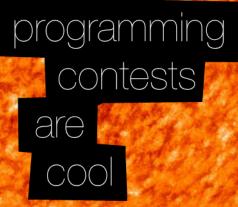


4	3	2	1	2	3	2	3	4
3	2	1		1	2	1	2	3
2	1	2	1	2	1		1	2
1		1	2	3	2	1	2	3
2	1	2	3	4	3	2	3	2
3	2	3	4	5	4	3	2	1
4	3	4	5	4	3	2	1	

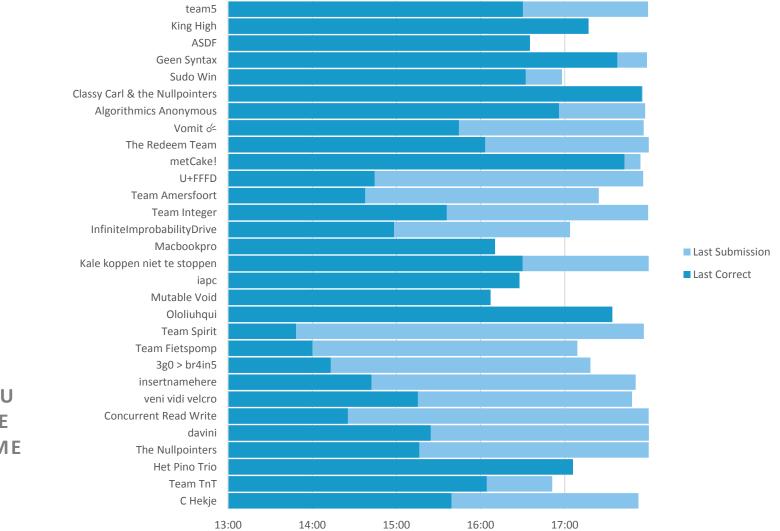


4	3	2	1	2	3	2	3	4
3	2	1		1	2	1	2	3
2	1	2	1	2	1		1	2
1		1	2	3	2	1	2	3
2	1	2	3	4	3	2	3	2
3	2	3	4	5	4	3	2	1
4	3	4	5	4	3	2	1	

Scores



#	AFFIL.	ТЕАМ	SCORE	A 😑	в 🔵	cO	D 🔾	Е 🦊	F	G 🔵	н●	I	ן 🔘 נ
1	r -	team5	9 142	1 (204 + 0)	2	2 (122 + 20)	1 (14 + 0)	1 (208 + 0)	1 (140 + 0)	1 (208 + 0)	1 (168 + 0)	2 (210 + 20)	2 (90 + 20)
2	۵	King High	8 98	0	2 (256 + 20)	1 (121 + 0)	1 (25 + 0)	1 (34 + 0)	2 (88 + 20)	3 (67 + 40)	3 (203 + 40)	2 (46 + 20)	1
3	14 -	ASDF	7 76	3 0	2 (174 + 20)	0	1 (137 + 0)	1 (12 + 0)	2 (52 + 20)	1 (215 + 0)	0	1 (22 + 0)	1 (111 + 0)
4	()	Geen Syntax	7 92	7 0	1 (164 + 0)	2	2 (132 + 20)	1 (15 + 0)	6 (106 + 100)	1 (36 + 0)	2	2 (37 + 20)	2 (277 + 20)
5	•	Sudo Win	6 64	90	0	3	1 (72 + 0)	1 (38 + 0)	3 (87 + 40)	2 (113 + 20)	0	1 (47 + 0)	2 (212 + 20)
6	۰ 🗖	Classy Carl & the Nullpointers	6 93	1	1	3	3 (148 + 40)	1 (34 + 0)	2 (93 + 20)	2 (294 + 20)	1	2 (84 + 20)	1 (181 + 0)
7	۰ 💳	Algorithmics Anonymous	5 55	71	0	0	4 (235 + 60)	1 (16 + 0)	2 (82 + 20)	1 (61 + 0)	0	3 (43 + 40)	1
8	TU/e	Vomit	4 35	2 1	0	0	1 (133 + 0)	1 (14 + 0)	4	1 (164 + 0)	0	1 (41 + 0)	1
9	•	The Redeem Team	4 42	31	1	0	1	1 (34 + 0)	6	2 (45 + 20)	0	3 (101 + 40)	1 (183 + 0)
10	TU/e	metCake!	4 83	1 0	2 (155 + 20)	0	6 (282 + 100)	2 (27 + 20)	0	7	0	5 (147 + 80)	0
11	TU/e	U+FFFD	3 21	90	o	o	2	1 (22 + 0)	5	3 (104 + 40)	o	2 (33 + 20)	0
12	•	Team Amersfoort	3 23	2 0	0	0	3	1 (61 + 0)	1 (97 + 0)	0	0	2 (54 + 20)	2
13	15 T	Team Integer	3 24	70	0	0	1	1 (18 + 0)	2	2 (155 + 20)	0	1 (54 + 0)	0
14	• =	Infinite Improbability Drive	3 29	50	2	0	0	1 (15 + 0)	1 (118 + 0)	1	0	5 (82 + 80)	0
15	(Macbookpro	3 35	PO	0	0	2 (190 + 20)	1 (21 + 0)	0	0	0	4 (68 + 60)	0
	()	Kale koppen niet te stoppen	3 37	۱o	0	0	3 (209 + 40)	1 (17 + 0)	2	0	0	3 (65 + 40)	0
17	9=	iapc	3 37	9 O	0	0	2 (207 + 20)	1 (16 + 0)	2	0	0	4 (76 + 60)	1
18	a 🗖	Mutable Void	3 45	0	0	0	0	1 (40 + 0)	0	5 (187 + 80)	0	2 (123 + 20)	0
19	• =	Ololiuhqui	3 80	3 0	o	0	2 (149 + 20)	1 (21 + 0)	0	0	0	18 (273 + 340)	0
20	15	Team Spirit	2 14	1	1	4	1	1 (48 + 0)	1	6	1	4 (33 + 60)	1
	14	Team Fietspomp	2 14	30	2	4	4	2 (28 + 20)	0	1	0	3 (60 + 40)	0
	a 🗖	3g0 > br4in5	2 16	2 0	0	3	0	1 (69 + 0)	3	0	0	2 (73 + 20)	0
23	9=	insertnamehere	2 16	10	0	0	0	1 (62 + 0)	0	0	0	1 (102 + 0)	1
24	15 T	veni vidi velcro	2 20	ιo	0	0	2	1 (26 + 0)	1	0	0	3 (135 + 40)	0
25	V	Concurrent Read Write	2 20	50	0	0	5 (85 + 80)	1 (40 + 0)	0	3	0	12	0
26	9=	davini	2 26	3 0	0	0	2	1 (59 + 0)	3	0	0	4 (144 + 60)	0
27	a 🗖	The Nullpointers	2 27	10	0	0	1	2 (35 + 20)	0	0	0	5 (136 + 80)	5
	٠	Het Pino Trio	2 45	2 0	0	o	0	1 (67 + 0)	o	0	o	8 (245 + 140)	0
29	• =	Team TnT	1 18	+0	0	0	0	1 (184 + 0)	3	0	0	7	0
30	ö	C Hekje	1 29	90	o	0	3	8 (159 + 140)	0	0	0	0	0
31	TU/e	phineas and ferb	0	0	0	0	1	8	0	0	0	2	1



WHEN YOU COULD'VE GONE HOME